



Approaches to distributed privacy protecting data mining

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Introduction

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- ⑥ Conflict Resolution: Inference Control
- ⑥ Inference Control Techniques
 - △ Controlled Release
 - △ Input/Output Perturbation
 - △ Query Restriction & Auditing
 - △ ...

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- ⑥ **Distributed** data mining – old and new challenges

Outline

- ⑥ What does “privacy” mean?
- ⑥ Perturbation techniques
- ⑥ Secure Multi-Party Computation (MPC)
- ⑥ Privacy by secure MPC
- ⑥ Conclusions

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 - △ Interval width for a confidence level [Agrawal, Srikant 2000]
 - △ Information theoretic approach [Agrawal, Aggarwal 2001]
 - △ Game theoretic approach [Kleinberg, Papadimitriou, Raghavan 2001]

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- ⑥ Privacy in secure multi-party computation

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- ⑥ Various techniques
 - △ randomized input distortion
 - △ output perturbation
 - △ k -anonymity [Sweeney '98]

Problems with Perturbations

- ⑥ Bias, precision & consistency
- ⑥ Can be computationally challenging
- ⑥ Outlier removal & “blurring” the data → detection of anomalies?
- ⑥ Combining multiple versions of data released for different purposes

Secure Multi-Party Computation

- ⑥ Introduced by Yao in 1982, inspired by “coin-flipping” (Blum) and “mental poker” (Shamir, Rivest, Adleman)
- ⑥ m parties P_1, \dots, P_m want to compute $f(x_1, \dots, x_m)$, where x_i is a private input of P_i , without revealing more than necessary ...

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- ⑥ ... i.e., simulation of a **trusted party!**
- ⑥ A very general and powerful tool, various models
- ⑥ Efficient completeness results: [Yao'86] (2-party), [GMW'87] (crypt.) and [BGW+CCD'88] (uncond.)

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- ⑥ Problems:
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 - △ Does it really solve the privacy problem?

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- ⑥ A different approach to MPC?

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- ⑥ Query restriction
 - △ query-set-size, query-set-overlap
 - △ query auditing
 - △ partitioning
- ⑥ Query auditing
 - △ efficient in simple cases
 - △ a NP-hard problem in general
[Kleinberg, Papadimitriou, Raghavan 2001]

Conclusions

- ⑥ “Privacy” means . . .
- ⑥ Various approaches, problem dependent
- ⑥ Probably no “the best” single solution
- ⑥ Still a lot of work to be done